

Lecture #15

• Last time:

- 1) $\mathbb{H}^f \circ \tau: M_2 \otimes M_1 \xrightarrow{\cong} M_1 \otimes M_2$ as $\mathcal{U}_q(\mathfrak{sl}_2)$ -mod
- 2) $\mathbb{H}^f_{12} \circ \mathbb{H}^f_{13} \circ \mathbb{H}^f_{23} = \mathbb{H}^f_{23} \circ \mathbb{H}^f_{13} \circ \mathbb{H}^f_{12}: M_1 \otimes M_2 \otimes M_3 \cong$

Given that f satisfies two difference equations:
 $f(\lambda, \mu) = \lambda f(\lambda, q^2 \mu) = \mu f(q^2 \lambda, \mu)$

Remark: a) There are many maps $f: \tilde{\Lambda} \times \tilde{\Lambda} \rightarrow \mathbb{K}^*$ satisfying these 2 difference eq-s (exercise: classify all).

b) In the simplest case of $M_1 = M_2 = L(1, +)$, it's easy to see that the map

$\mathbb{H}: L(1, +)^{\otimes 2} \cong$ in the basis $\{v_0 \otimes v_0, v_0 \otimes v_1, v_1 \otimes v_0, v_1 \otimes v_1\}$ is given by

$$\mathbb{H} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & q^{-1}q & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

c) Since the weights of $L(1, +)$ are $q^{\pm 1}$ (as $Kv_0 = q \cdot v_0, Kv_1 = q^{-1} \cdot v_1$), we see that f is uniquely determined by $\lambda := f(q, q)$, since $f(q, q^{-1}) = q \cdot f(q, q) = f(q^{-1}, q)$ and $f(q^{-1}, q^{-1}) = q^{-1} \cdot f(q^{-1}, q)$ so that $f(q, q^{-1}) = q \cdot \lambda, f(q^{-1}, q) = q \cdot \lambda, f(q^{-1}, q^{-1}) = \lambda$. Thus:

$$\mathbb{H}^f = \begin{pmatrix} \lambda & 0 & 0 & 0 \\ 0 & \lambda q & 0 & 0 \\ 0 & \lambda q(q^{-1}-q) & \lambda q & 0 \\ 0 & 0 & 0 & \lambda \end{pmatrix}$$

The following result allows to relate $\mathcal{U}_q(\mathfrak{sl}_2)$ -modf.d. to tensor categories:

Proposition 1: Let M_1, M_2, M_3 be fin. dim. $\mathcal{U}_q(\mathfrak{sl}_2)$ -modules. We also assume

that f in addition is multiplicative in each argument:

$$f(\lambda, \mu\nu) = f(\lambda, \mu)f(\lambda, \nu), \quad f(\lambda\mu, \nu) = f(\lambda, \nu)f(\mu, \nu)$$

for all weights λ, μ, ν of these modules. Then:

$$M_1 \otimes (M_2 \otimes M_3) \xrightarrow{\text{id} \otimes \check{R}} M_1 \otimes (M_3 \otimes M_2) \xrightarrow{\text{can}} (M_1 \otimes M_3) \otimes M_2 \xrightarrow{\check{R} \otimes \text{id}} (M_3 \otimes M_1) \otimes M_2$$

$$\xrightarrow{\text{can}} (M_1 \otimes M_2) \otimes M_3 \xrightarrow{\check{R}} M_3 \otimes (M_1 \otimes M_2) \xrightarrow{\text{can}} (M_3 \otimes M_1) \otimes M_2$$

and

$$(M_1 \otimes M_2) \otimes M_3 \xrightarrow{\check{R} \otimes \text{id}} (M_2 \otimes M_1) \otimes M_3 \xrightarrow{\text{can}} M_2 \otimes (M_1 \otimes M_3) \xrightarrow{\text{id} \otimes \check{R}} M_2 \otimes (M_3 \otimes M_1)$$

$$\xrightarrow{\text{can}} (M_1 \otimes M_2) \otimes M_3 \xrightarrow{\check{R}} (M_2 \otimes M_3) \otimes M_1 \xrightarrow{\text{can}} (M_2 \otimes M_3) \otimes M_1$$

where $\text{can}: (A \otimes B) \otimes C \xrightarrow{\cong} A \otimes (B \otimes C)$ is the canonical associativity map ($(a \otimes b) \otimes c \mapsto a \otimes (b \otimes c)$)

$\check{R} := \mathbb{H}^f \circ \tau: B \otimes A \xrightarrow{\cong} A \otimes B$ is the isom. of $\mathcal{U}_q(\mathfrak{sl}_2)$ -modf.d.

Lecture # 15

Proof of Proposition 1

The "upper" composition in the first diagram reads:

$$\mathbb{H}_{12} \circ \tilde{f}_{12} \circ \tau_{12} \circ \mathbb{H}_{23} \circ \tilde{f}_{23} \circ \tau_{23} = \mathbb{H}_{12} \circ \tilde{f}_{12} \circ \mathbb{H}_{13} \circ \tau_{12} \circ \tilde{f}_{23} \circ \tau_{23} = \mathbb{H}_{12} \circ \tilde{f}_{12} \circ \mathbb{H}_{13} \circ \tilde{f}_{13} \circ \tau_{12} \circ \tau_{23}$$

But, due to Lemma 1 of Lecture 14, we have:

$$\tilde{f}_{12} \circ \mathbb{H}_{13} = \mathbb{H}' \circ \tilde{f}_{12}$$

Finally, by analogy with the proof of Lemma 3 of Lecture 14, we have:

$$\mathbb{H}_{12} \circ \mathbb{H}' = (\mathbb{1} \otimes \Delta)(\mathbb{H})$$

Exercise: Prove this.

Thus, the upper composition of the 1st diagram reads:

$$(\mathbb{1} \otimes \Delta)(\mathbb{H}) \circ (\tilde{f}_{12} \circ \tilde{f}_{13}) \circ (\tau_{12} \circ \tau_{23})$$

On the other hand, evaluating the lower composition of the 1st diagram on the tensor product of weight vectors $m_1 \otimes m_2 \otimes m_3 \in (M_1)_\lambda \otimes (M_2)_\mu \otimes (M_3)_\nu$

we get:

$$m_1 \otimes (m_2 \otimes m_3) \xrightarrow{\text{can}} \underbrace{(m_1 \otimes m_2)}_{(M_1 \otimes M_2)_{\lambda+\mu}} \otimes m_3 \xrightarrow{\tau} m_3 \otimes (m_1 \otimes m_2) \xrightarrow{\mathbb{H}^f} f(\nu, \lambda, \mu) \circ (\mathbb{1} \otimes \Delta)(\mathbb{H})(m_3 \otimes (m_1 \otimes m_2))$$

As $\tau_{12} \circ \tau_{23}: m_1 \otimes m_2 \otimes m_3 \mapsto m_3 \otimes m_1 \otimes m_2$

$\tilde{f}_{12} \circ \tilde{f}_{13}: m_3 \otimes m_1 \otimes m_2 \mapsto f(\nu, \lambda) \cdot f(\nu, \mu) \cdot m_3 \otimes m_1 \otimes m_2$

we conclude that the 1st diagram is indeed commutative due to

$$f(\nu, \lambda, \mu) = f(\nu, \lambda) f(\nu, \mu)$$

Likewise, the 2nd diagram is commutative due to $f(\lambda, \mu, \nu) = f(\lambda, \nu) f(\mu, \nu)$

Def: A finite dimensional q -module is called of type 1 if all eigenvalues of K are powers of q (no minus sign!)

Let $\Lambda = \{q^s \mid s \in \mathbb{Z}\} \subseteq \tilde{\Lambda} = \{\pm q^s \mid s \in \mathbb{Z}\}$

Exercise: a) Classify all $f: \Lambda \times \Lambda \rightarrow \mathbb{k}^*$ satisfying 4 conditions
 $f(\lambda, \mu) = \lambda f(\lambda, q^2 \mu) = \mu f(q^2 \lambda, \mu)$, $f(\lambda, \mu, \nu) = f(\lambda, \nu) f(\mu, \nu)$, $f(\lambda, \mu, \nu) = f(\lambda, \mu) f(\lambda, \nu)$
b) Classify all $f: \tilde{\Lambda} \times \tilde{\Lambda} \rightarrow \mathbb{k}^*$ satisfying the above 4 conditions

Hint to Exercise

a) Combine $f(q^r, q^s) = f(q, q)^{rs}$ with $f(q, q)^2 = q^{-1} f(q, 1)$ to deduce $f(q, q) = (q^{-1/2})^{rs}$ given \mathbb{k} contains a square root of q .

b) Combine $f(-1, 1)^2 = f(-1, 1)$ with $f(1, q) = f(-1, q)^2 = -f(-1, 1)$ to get a contradiction.

Remark: The absence of $f: \mathbb{A} \times \mathbb{A} \rightarrow \mathbb{k}^*$ satisfying these 4 conditions explains why people usually consider only type 1 fundam. $U_q(\mathfrak{sl}_2)$ -modules in the relation to provided tensor categories. The latter are the categories \mathcal{C} equipped with functor $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ together with

$$(X, Y) \mapsto X \otimes Y$$

functorial isomorphisms (called associativity/commutativity constraints)

$$(X \otimes Y) \otimes Z \cong X \otimes (Y \otimes Z) \quad \text{and} \quad X \otimes Y \cong Y \otimes X.$$

subject to some compatibilities, including the two diagrams from Prop 1, called the hexagonal diagrams.

For the rest of today (and next lecture) we will try to answer:

Question: What is the exact relation between the quantum groups $U_q(\mathfrak{sl}_2)$, $Sl_q(\mathfrak{sl}_2)$ that we discussed in the first 2 weeks and $U_q(\mathfrak{sl}_2)$ that we have been discussing over the last 2 weeks?

To get an answer, we shall first understand the classical ($q=1$) case.

To this end, we start with the following definition.

Def: a) Given two bialgebras A and B , a bilinear form $\langle \cdot, \cdot \rangle: A \times B \rightarrow \mathbb{k}$ is called a bialgebra pairing if the following holds (using Sweedler's notation):

$$\langle a, b_1 b_2 \rangle = \sum_{(a)} \langle a', b_1 \rangle \langle a'', b_2 \rangle, \quad \langle a_1 a_2, b \rangle = \sum_{(b)} \langle a_1, b' \rangle \langle a_2, b'' \rangle$$

$$\langle a, 1_B \rangle = \varepsilon_A(a), \quad \langle 1_A, b \rangle = \varepsilon(b)$$

b) If in addition A, B are Hopf algebras, then such $\langle \cdot, \cdot \rangle$ is a Hopf pairing if also

$$\langle S_A(a), b \rangle = \langle a, S_B^{-1}(b) \rangle$$

Remark: The above definition differs slightly from (Kassel, Def V.7.17).
To get his setup, take $A = \mathcal{U}$, $B = H^{\text{cop}}$.

The following simple result explains the origin of this definition

Lemma 1: a) Given bialgebras A, B and a bilinear pairing $\langle, \rangle: A \times B \rightarrow k$
 \langle, \rangle is a bialgebra pairing iff the natural linear maps

$$\varphi: A^{\text{op}} \rightarrow B^* \quad \varphi(a)(b) = \langle a, b \rangle$$

$$\psi: B \rightarrow A^* \quad \psi(b)(a) = \langle a, b \rangle$$
 are algebra morphisms.

b) If $\dim(B) < \infty$, then \langle, \rangle -bialgebra pairing iff above
 $\varphi: A^{\text{op}} \rightarrow B^*$ is a bialgebra morphism.

a) Let's show that $\langle a_1 a_2, b \rangle = \sum_{(b)} \langle a_1, b' \rangle \langle a_2, b'' \rangle$ is equivalent to $\varphi(a_1 a_2) = \varphi(a_2) \varphi(a_1)$

Indeed:

$$\varphi(a_1 a_2)(b) = \langle a_1 a_2, b \rangle \quad \text{while} \quad (\varphi(a_2) \varphi(a_1))(b) = \sum_{(b)} \langle a_2, b' \rangle \cdot \langle a_1, b'' \rangle.$$

Also: $\varphi(1_A)(b) = \langle 1_A, b \rangle = \varepsilon_B(b)$ implies $\varphi(1_A) = 1_{B^*}$ and vice-versa.

Likewise, $\langle a, b_1 b_2 \rangle = \sum_{(a)} \langle a', b_1 \rangle \langle a'', b_2 \rangle$ is equivalent to
 $\psi(b_1 b_2) = \psi(b_1) \psi(b_2)$, while $\langle a, 1_B \rangle = \varepsilon_A(a)$ is equivalent to $\psi(1_B) = 1_{A^*}$.

b) It suffices to check that $\begin{cases} \langle a, b_1 b_2 \rangle = \sum_{(a)} \langle a', b_1 \rangle \langle a'', b_2 \rangle \\ \langle a, 1_B \rangle = \varepsilon_A(a) \end{cases}$ are equivalent

to φ being a coalgebra morphism. The first follows from

$$\Delta(\varphi(a))(b \otimes c) = \varphi(a)(bc) = \langle a, bc \rangle = \sum_{(a)} \langle a', b \rangle \langle a'', c \rangle = (\varphi \otimes \varphi)(\Delta(a))(b \otimes c)$$

while the 2nd is also clear

Next time: construct a bialgebra pairing $M(2) \times \mathcal{U}(\mathfrak{sl}_2) \rightarrow k$ which gives rise to a Hopf pairing $SL(2) \times \mathcal{U}(\mathfrak{sl}_2) \rightarrow k$. Both will naturally generalize to q -version.

Remark: A geometric way to think about the above pairing btw SL_2 and $U(\mathfrak{sl}_2)$ is by realizing $U(\mathfrak{sl}_2)$ = left/right-invariant differential operators on $SL_2 = \text{Spec}(SL_2)$ and evaluating the application of this operator to f of SL_2 at $1 \in SL_2$.