

Lecture #25

Last time we constructed a pairing $U_q^{\leq} \times U_q^{\geq} \rightarrow \mathbb{k}$ s.t.

$$(y, x_1 x_2) = (\Delta(y), x_2 \otimes x_1) \quad \text{and} \quad (y_1 y_2, x) = (y_1 \otimes y_2, \Delta(x))$$

which becomes a bialgebra pairing in the sense of Lecture 15 after the swap of factors.

Similarly to our proof of [Lecture 24, Lemma 6], the above f-laws allow one to compute (y, x) iteratively by using Q-greedy reasoning and f-laws for coproduct.

To this end recall [Lecture 18, Lemma 1]:

$$\Delta((U_q^+)_{\mu}) \subseteq \bigoplus_{0 \leq \nu \leq \mu} (U_q^+)_{\mu-\nu} K_{\nu} \otimes (U_q^+)_{\nu}$$

As $(U_q^+)_{\alpha_i} = \mathbb{k} \cdot E_i$ for any simple root α_i , one can introduce linear maps

$$\tau_i, \tau'_i: U_q^+ \rightarrow U_q^+ \quad \forall i \in I$$

via the following equality:

$$\Delta(x) = x \otimes 1 + \sum_{i \in I} \tau_i(x) K_i \otimes E_i + \dots + \sum_{i \in I} E_i K_{\mu-\alpha_i} \otimes \tau'_i(x) + K_{\mu} \otimes x \quad \forall x \in (U_q^+)_{\mu}$$

The reasoning behind this definition is precisely part (b) of the next lemma:

Lemma 1: a) For $x \in (U_q^+)_{\mu}$, $x' \in (U_q^+)_{\mu'}$, we have

$$\begin{aligned} \tau_i(xx') &= x \cdot \tau_i(x') + q^{(\alpha_i, \mu')} \tau_i(x) \cdot x' \\ \tau'_i(xx') &= q^{(\alpha_i, \mu)} x \cdot \tau'_i(x') + \tau'_i(x) \cdot x'. \end{aligned}$$

$$b) (F_i y, x) = (F_i, E_i) \cdot (y, \tau'_i(x)) \quad , \quad (y F_i, x) = (F_i, E_i) \cdot (y, \tau_i(x))$$

$$c) \tau'_i = \sigma \tau_i \sigma \quad \text{with } \sigma: U_q \ni \text{ being the antiautomorphism mapping } \begin{matrix} E_i \mapsto E_i, F_i \mapsto F_i, \\ K_i^{\pm 1} \mapsto K_i^{\mp 1} \end{matrix}$$

$$\Rightarrow a) \text{ If } \Delta(x) = x \otimes 1 + \sum_{i \in I} \tau_i(x) K_i \otimes E_i + \dots \quad , \quad \Delta(x') = x' \otimes 1 + \sum_{i \in I} \tau_i(x') K_i \otimes E_i + \dots \quad ,$$

$$\text{then } \Delta(xx') = \Delta(x) \Delta(x') = xx' \otimes 1 + \sum_{i \in I} \underbrace{(x \tau_i(x') K_i + \tau_i(x) K_i x')}_{(x \tau_i(x') + q^{(\alpha_i, \mu')} \tau_i(x) x') K_i} \otimes E_i + \dots$$

which proves the 1st f-law in a). The 2nd f-law is established similarly.

$$b) (F_i y, x) = (F_i \otimes y, \dots + \sum_j E_j K_{\mu-\alpha_j} \otimes \tau'_i(x) + K_{\mu} \otimes x) \xrightarrow[\text{reasons}]{\text{degree}} \underbrace{(F_i, E_i K_{\mu-\alpha_i})}_{=(F_i, E_i)} \cdot (y, \tau'_i(x))$$

which proves the 1st f-law in (b). The 2nd f-law is proved similarly.

[Exercise: Prove part c)

The above construction can be likewise applied to U_q^- . To this end, we define

$$\tau_i, \tau_i' : U_q^- \rightarrow U_q^- \quad \forall i \in I$$

via the following equality:

$$\Delta(y) = y \otimes K_{-\mu} + \sum_{i \in I} \tau_i(y) \otimes F_i K_{-(\mu-d_i)^+} + \sum_{i \in I} F_i \otimes \tau_i'(y) K_i + 1 \otimes y \quad \forall y \in (U_q^-)_{\mu}$$

The following is the analogue of Lemma 1:

Lemma 2: a) For $y \in (U_q^-)_{-\mu}$, $y' \in (U_q^-)_{-\mu'}$, we have

$$\tau_i(y y') = q^{(x_i, \mu')} y \tau_i(y') + \tau_i(y) \cdot y'$$

$$\tau_i'(y y') = y \tau_i'(y') + q^{(x_i, \mu)} \tau_i'(y) y'$$

$$b) (y, E_i x) = (F_i, E_i) \cdot (\tau_i(y), x) \quad , \quad (y, x E_i) = (F_i, E_i) \cdot (\tau_i'(y), x)$$

$$c) \tau_i' = \delta \tau_i \delta$$

d) We have $\tau_i(y) = \omega \tau_i' \omega(y)$ and $\tau_i'(y) = \omega \tau_i \omega(y) \quad \forall y \in U_q^-$,
where $\omega : U_q \rightarrow U_q$ is the Cartan involution mapping $E_i \mapsto F_i, F_i \mapsto E_i, K_i^{\pm 1} \mapsto K_i^{\mp 1}$

a) - c) are proved entirely analogously to Lemma 1.

The equalities of d) hold for $y = F_j \quad \forall j \in I$, hence, it suffices to check that if d) holds for y_1, y_2 , then it also holds for $y_1 \cdot y_2$. This is straightforward:

$$\begin{aligned} \omega \tau_i' \omega(y_1 y_2) &= \omega \tau_i'(\omega(y_1) \cdot \omega(y_2)) \stackrel{\text{Lemma 1}}{=} \omega(q^{(x_i, \mu')} \omega(y_1) \cdot \tau_i' \omega(y_2) + \tau_i' \omega(y_1) \cdot \omega(y_2)) \\ &= q^{(x_i, \mu')} \cdot \underbrace{\omega^2(y_1)}_{=y_1} \cdot \underbrace{\omega \tau_i' \omega(y_2)}_{=\tau_i'(y_2)} + \underbrace{\omega \tau_i' \omega(y_1)}_{=\tau_i'(y_1)} \cdot \underbrace{\omega^2(y_2)}_{=y_2} \stackrel{\text{Lemma 2a)}}{=} \tau_i(y_1 y_2) \end{aligned}$$

The 2nd of the in d) is checked similarly

$$\text{Lemma 3: } (y, x) = (\omega(x), \omega(y)) \quad \forall y \in U_q^-, x \in U_q^+$$

Can assume $x \in (U_q^+)_{\mu}$, $y \in (U_q^-)_{-\mu}$. Moreover, suffices to treat $y = F_i y', y' \in (U_q^-)_{-\mu+d_i}$:

$$\begin{aligned} (F_i y', x) &\stackrel{\text{Lemma 1}}{=} (F_i, E_i) \cdot (y', \tau_i'(x)) \stackrel{\text{Assumption of induction}}{=} (F_i, E_i) \cdot (\omega \tau_i'(x), \omega(y')) \stackrel{\text{Lemma 2d)}}{=} \\ &(F_i, E_i) \cdot (\tau_i \omega(x), \omega(y')) \stackrel{\text{Lemma 2}}{=} (\omega(x), \underbrace{E_i}_{=\omega(F_i)} \omega(y')) = (\omega(x), \omega(y')) \end{aligned}$$

Thus, the proof proceeds by induction on the height of $\mu \in Q$

Exercise: Prove that $(S(y), S(x)) = (y, x) \quad \forall y \in U_q^-, x \in U_q^+$, where S -antipode.

More importantly, the maps τ_i, τ'_i also arise when commuting elements of U_q^- with E_i , or el-s of U_q^+ with F_i , as shown in the next result:

Lemma 4: For $i \in I$, $y \in (U_q^-)_{-\mu}$, $x \in (U_q^+)_{\mu}$, we have:

$$E_i y - y E_i = \frac{1}{q_i - q_i^{-1}} (K_i \tau_i(y) - \tau'_i(y) K_i^{-1})$$

$$F_i x - x F_i = \frac{-1}{q_i - q_i^{-1}} (\tau_i(x) K_i - K_i^{-1} \tau'_i(x))$$

► First, we prove the first equality. It's obvious for $y = F_j$. It thus remains to check that it holds for $y_1 \cdot y_2$ if it holds for y_1 & y_2 :

$$\begin{aligned} [E_i, y_1 \cdot y_2] &= [E_i, y_1] \cdot y_2 + y_1 \cdot [E_i, y_2] = \frac{1}{q_i - q_i^{-1}} (K_i \tau_i(y_1) y_2 - \tau'_i(y_1) K_i^{-1} y_2) + \\ &+ \frac{1}{q_i - q_i^{-1}} (y_1 K_i \tau_i(y_2) - y_1 \tau'_i(y_2) K_i^{-1}) = \frac{1}{q_i - q_i^{-1}} K_i \left(\underbrace{\tau_i(y_1) y_2 + q^{(\alpha_i, w^{\vee} y_1)} y_1 \cdot \tau_i(y_2)}_{= \tau_i(y_1 y_2)} \right) - \\ &- \frac{1}{q_i - q_i^{-1}} \left(\underbrace{\tau'_i(y_1) \cdot y_2 + q^{(\alpha_i, w^{\vee} y_2)} y_1 \tau'_i(y_2)}_{= \tau'_i(y_1 y_2)} \right) K_i^{-1} \end{aligned}$$

To deduce the second equality in the lemma, apply Cartan involution w to the 1st one with $y = w(x)$, and use Lemma 4b). ■

The above result implies also:

Lemma 5: For any $i \in I$, $\lambda \in Q$, $x \in (U_q^+)_{\mu}$, $y \in (U_q^-)_{-\nu}$, we have:

$$\text{ad}(E_i)(y K_{\lambda} x) = y K_{\lambda} (q^{-(\lambda, \alpha_i)} E_i x - q^{(\mu - \nu, \alpha_i)} x E_i) + \frac{(q^{-(\nu - \alpha_i, \alpha_i)} \tau_i(y) K_{\lambda + \alpha_i} - \tau'_i(y) K_{\lambda - \alpha_i}) x}{q_i - q_i^{-1}}$$

$$\text{ad}(F_i)(y K_{\lambda} x) = q^{-(\mu, \alpha_i)} (F_i y - q^{-(\lambda, \alpha_i)} y F_i) K_{\lambda + \alpha_i} x + \frac{y (q^{-(\mu - \alpha_i, \alpha_i)} K_{\lambda} \tau'_i(x) - q^{-2(\mu - \alpha_i, \alpha_i)} K_{\lambda + 2\alpha_i} \tau_i(x))}{q_i - q_i^{-1}}$$

Exercise: Prove this lemma.

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We conclude with the following main result for today:

Proposition 1: Assuming "tQ-condition" (see Lecture 23), the restriction $(\cdot, \cdot) : (U_q^-)_{-\mu} \times (U_q^+)_{\mu} \rightarrow k$ is non-degenerate $\forall \mu \neq 0$.

As $\dim(U_q^-)_{-\mu} = \dim(U_q^+)_{\mu}$, due to Cartan involution, it suffices to show $(y, (U_q^+)_{\mu}) = 0 \Rightarrow y = 0$. The proof is by induction on the height of μ . Applying Lemma 2 to $(y, E_i x')$ and $(y, x' E_i)$ for any $i \in I$ and $x' \in (U_q^+)_{\mu - \alpha_i}$ and using the assumption of induction, we find $r_i(y) = 0 = r_i'(y) \forall i$.

But then, $[y, E_i] = 0 \forall i \in I$, due to Lemma 4. Applying ω to this equality, we see that $x := \omega(y) \in (U_q^+)_{\mu}$ commutes with all F_i .

Claim: If $x \in (U_q^+)_{\mu}$ ($\mu > 0$) commutes with all $\{F_i\}_{i \in I}$, then $x = 0$

For any $\lambda \in P$, consider the Verma module $M(\lambda)$ with the highest weight vector v_{λ} . For degree reasons: $x(v_{\lambda}) = 0$. But as $M(\lambda)$ is generated from v_{λ} by $\{F_j\}$ which commute with x , we conclude that $x|_{M(\lambda)} = 0$. Hence, x acts trivially on all simple finite-dimensional modules.

But by [Lecture 22, Theorem 2], under "tQ-condition" any finite-dimen. $U_q(\mathfrak{g})$ -module is $\cong \bigoplus$ simple f.d., hence, x acts trivially on all f.d. mod. The latter implies the stated $x = 0$, due to [Huk 4, Problem 4], cf. our proof of [Lecture 23, Lemma 2b)]. □

This establishes the above Claim, thus completing the proof of Prop. Combining this result and Lemma 3, we obtain:

Corollary 1: The bilinear form $(\cdot, \cdot) : U_q^+ \times U_q^+ \rightarrow k$ given by $(x_1, x_2) = (\omega(x_1), x_2)$ is symmetric and non-degenerate.