

## Lecture #42

Goal: Braided bialgebras, cobraided bialgebras, and the FRT construction.

Last time we discussed braided tensor categories and at the very end asked:

Q: Given a bialgebra  $H$ , which conditions on  $H$  are equivalent to  $H\text{-Mod}$  being braided?

To answer this, we start with the following key definition:

Def: a) A bialgebra  $H$  is quasi-cocommutative if  $\exists R \in H \otimes H$  satisfying:

$$\boxed{R\text{-invertible} \quad \text{and} \quad \Delta^{\otimes}(x) = R \Delta(x) R^{-1} \quad \forall x \in H}$$

Such  $R$  is called a universal R-matrix

b) A quasi-cocommutative bialgebra  $H$  with universal R-matrix  $R$  is braided if:

$$\boxed{(\Delta \otimes \text{id})(R) = R_{13} R_{23} \quad \text{and} \quad (\text{id} \otimes \Delta)(R) = R_{13} R_{12}}$$

Finally, if  $H$  had the antipode  $S$  (thus being a Hopf algebra), then we say that  $H$  is a quasi-cocommutative Hopf algebra, respectively braided Hopf alg, if the above holds.

Proposition 1: If  $(H, \mu, \eta, \Delta, \varepsilon, R)$  is a braided bialgebra, then:

a)  $R_{12} R_{13} R_{23} = R_{23} R_{13} R_{12}$

b)  $(\varepsilon \otimes \text{id})(R) = 1 = (\text{id} \otimes \varepsilon)(R)$

If furthermore  $H$  is a Hopf algebra with antipode  $S$ , then:

c)  $(S \otimes \text{id})(R) = R^{-1} = (\text{id} \otimes S^{-1})(R)$

d)  $(S \otimes S)(R) = R^{-1}$

► a)  $R_{12} R_{13} R_{23} = R_{12} (\Delta \otimes \text{id})(R) = (\Delta^{\otimes} \otimes \text{id})(R) R_{12} \stackrel{\text{erc } (\varepsilon \otimes \text{id}) \Delta = \text{id}}{=} (\varepsilon \otimes \text{id})(\underbrace{(\Delta \otimes \text{id})(R)}_{R_{13} R_{23}}) \cdot R_{12} = R_{23} R_{13} R_{12}$

b)  $R \stackrel{\text{erc } (\varepsilon \otimes \text{id} \otimes \text{id})(\Delta \otimes \text{id})(R) = (\varepsilon \otimes \text{id} \otimes \text{id})(R_{13} R_{23}) = (\varepsilon \otimes \text{id})(R) \cdot R}{=} \stackrel{\text{R-invertible}}{=} (\varepsilon \otimes \text{id})(R) = 1$

Using  $(\text{id} \otimes \varepsilon) \Delta = \text{id}$  instead one gets  $(\text{id} \otimes \varepsilon)(R) = 1$

c)  $\underbrace{(\mu \otimes \text{id})(S \otimes \text{id} \otimes \text{id})(\Delta \otimes \text{id})(R)}_{(S \otimes \text{id})(R) \cdot R} = \underbrace{(\varepsilon \otimes \text{id})(R)}_{\text{erc } \mu(S \otimes \text{id}) \Delta = \varepsilon} \stackrel{b)}{=} 1 \Rightarrow (S \otimes \text{id})(R) = R^{-1}$

The other equality is similar.

d)  $(S \otimes S)(R) = (\text{id} \otimes S)(S \otimes \text{id})(R) \stackrel{c)}{=} (\text{id} \otimes S)(\text{id} \otimes S^{-1})(R) = (\text{id} \otimes \text{id})(R) = R$

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Given any  $V, W \in H\text{-Mod}$ , where  $(H, \mu_H, \Delta, \varepsilon, R)$  - braided bialgebra, define

$$\boxed{C_{V,W}^R : V \otimes W \xrightarrow{\sim} W \otimes V \\ v \otimes w \mapsto \tau_{V,W}(R(v \otimes w))}$$

Proposition 2: The above maps  $C_{V,W}^R$  satisfy:

- a)  $C_{V,W}^R$  - isomorphism of  $H$ -modules
- b)  $\forall U, V, W \in H\text{-Mod}$ :

$$C_{U \otimes V, W}^R = (C_{U,W}^R \otimes \text{id}_V)(\text{id}_U \otimes C_{V,W}^R)$$

$$C_{U, V \otimes W}^R = (\text{id}_U \otimes C_{V,W}^R)(C_{U,V}^R \otimes \text{id}_W)$$

- c)  $\forall U, V, W \in H\text{-Mod}$ :

$$(C_{V,W}^R \otimes \text{id}_U)(\text{id}_V \otimes C_{U,W}^R)(C_{U,V}^R \otimes \text{id}_W) = (\text{id}_W \otimes C_{U,V}^R)(C_{U,W}^R \otimes \text{id}_V)(\text{id}_U \otimes C_{V,W}^R)$$

This is an immediate consequence of the defining relations of  $R$  (combined with  $\tau$ )

$$a) C_{V,W}^R(x(v \otimes w)) = \tau_{V,W} R \Delta(x)(v \otimes w) = \tau_{V,W} \Delta^\text{op}(x) R(v \otimes w) = \Delta(x) \tau_{V,W} R(v \otimes w) \\ x (C_{V,W}^R(v \otimes w))$$

b) It will be useful to write one of the braiding cond's on  $R$  as:

$$\sum_{i,j} s_i s_j \otimes t_j \otimes t_i = \sum_{i, (t_i)} s_i \otimes t'_i \otimes t''_i \quad \text{where } R = \sum_i s_i \otimes t_i \in H \otimes H.$$

Then, the 2nd equality of b) follows:

$$(\text{id}_V \otimes C_{U,W}^R)(C_{U,V}^R \otimes \text{id}_W)(U \otimes v \otimes w) = (\text{id}_V \otimes C_{U,W}^R)(\sum_i t_i v \otimes s_i u \otimes w) = \\ = \sum_{i,j} t_i v \otimes t_j w \otimes s_j s_i u \stackrel{\substack{\text{above} \\ \text{equality}}}{=} \sum_{i, (t_i)} t'_i v \otimes t''_i w \otimes s_i u = \underbrace{\sum_i t_i (v \otimes w) \otimes s_i u}_{C_{U,V \otimes W}^R(U \otimes v \otimes w)}$$

The 1st equality of b) follows likewise

c) Follows from Prop 1a), which can be interpreted as:

$$\sum_{i,j,k} s_k s_j \otimes t_k s_i \otimes t_j t_i = \sum_{i,j,k} s_j s_i \otimes s_k t_i \otimes t_k t_j$$

Indeed:

$$\begin{aligned} \text{LHS}(u \otimes v \otimes w) &= (C_{V,W}^R \otimes \text{id}_U)(\text{id}_V \otimes C_{U,W}^R)(\sum_i t_i v \otimes s_i u \otimes w) = (C_{V,W}^R \otimes \text{id}_U)(\sum_{i,j} t_i v \otimes t_j w \otimes s_j s_i u) \\ &\quad \left( \begin{array}{l} \text{above} \\ \text{equality} \end{array} \right) \\ &= \sum_{i,j,k} t_k t_j w \otimes s_k t_i v \otimes s_j s_i u \end{aligned}$$

$$\begin{aligned} \text{RHS}(u \otimes v \otimes w) &= (\text{id}_W \otimes C_{U,V}^R)(C_{U,W}^R \otimes \text{id}_V)(\sum_i u \otimes t_i w \otimes s_i v) = (\text{id}_W \otimes C_{U,V}^R)(\sum_{i,j} t_j t_i w \otimes s_j s_i u) \\ &= \sum_{i,j,k} t_j t_i w \otimes t_k s_i v \otimes s_k s_j u \end{aligned}$$

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Def: Given a vector space  $V$ , an invertible endomorphism  $c \in \text{Aut}(V \otimes V)$  is called an R-matrix if it satisfies the Yang-Baxter equation:  $c_{12} c_{23} c_{12} = c_{23} c_{12} c_{23}$  (viewed as equality in  $\text{End}(V \otimes V \otimes V)$  with  $c_{12} = c \otimes \text{id}_V$ ,  $c_{23} = \text{id}_V \otimes c$ )

Taking  $U=V=W$  in Proposition 2.c), we thus get:

Corollary 1: If  $(H, \mu, \gamma, \Delta, \varepsilon, R)$ -braided bialgebra, then for any  $V \in H\text{-Mod}$ ,  $c_{V,V}^R \in \text{End}(V \otimes V)$  is an R-matrix (in above sense)

Now we can finally answer the question from the very beginning:

Proposition 3:  $H\text{-Mod}$  is braided  $\Leftrightarrow H$  is braided



The braiding on the category  $H\text{-Mod}$  is precisely given by  $\{c_{V,W}^R \mid V, W \in H\text{-Mod}\}$

The conditions of a braiding on a category then follow from Prop. 2.



Assume we have a braiding on the category  $H\text{-Mod}$ . First, we produce  $R$ :

$$R := \tau_{H,H}(c_{H,H}(1 \otimes 1)) \quad \text{where } c_{H,H}: H \otimes H \rightarrow H \otimes H$$

Then: 1) naturality of  $c_{..}$  will imply R-universal R-matrix  
2) hexagon axiom will imply the braiding equalities

But first we note that the above construction is indeed "inverse" to  $\Leftarrow$ :

for any  $V, W \in H\text{-Mod}$ , and any  $v \in V, w \in W$ , consider morphisms  $H \xrightarrow{\varphi_v} V$   $H \xrightarrow{\varphi_w} W$   
s.t.  $\varphi_v(h) = hv$ ,  $\varphi_w(h) = hw$ . Then naturality of braiding implies:

$$\begin{array}{ccc} H \otimes H & \xrightarrow{c_{H,H}} & H \otimes H \\ \varphi_v \otimes \varphi_w \downarrow & & \downarrow \varphi_w \otimes \varphi_v \\ V \otimes W & \xrightarrow{c_{V,W}} & W \otimes V \end{array}$$

Tracing the images of  $1 \otimes 1 \in H \otimes H$ , we get:  $c_{V,W}(v \otimes w) = \tau_{V,W}(R(v \otimes w)) = c_{V,W}^R(v \otimes w)$

Back to 1) from above. Note  $c_{H,H}$  is an  $H$ -module morphism  $\Rightarrow x \circ c_{H,H} = c_{H,H} \circ x$

$$\forall x \in H \implies \Delta(x)(c_{H,H}(1 \otimes 1)) = c_{H,H}(\Delta(x)(1 \otimes 1)) \stackrel{\text{above}}{=} \tau_{H,H}^R(R \Delta(x)) \Rightarrow \Delta^{\text{op}}(x)R = R \Delta(x)$$

As per 2) - direct chase of diagram

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By the above result, a braiding on the category  $H\text{-Mod}$  is equivalent to a braiding on the bialgebra  $H$ . There is a decal notion of cobrading on  $H$ ! To motivate the definition, let's view  $H \otimes H$  as  $\text{Hom}_{\mathbb{K}}(\mathbb{K}, H \otimes H)$ , so that  $R \in H \otimes H$  from before is really a linear map  $\mathbb{K} \xrightarrow{\quad R \quad} H \otimes H$ . Thus, following our lectures 1-3, the dual object should be a linear map  $H \otimes H \rightarrow \mathbb{K}$ .

Def: a) A bialgebra  $H$  is quasi-cocommutative if there is a linear form  $r \in (H \otimes H)^*$

- $\exists \bar{r} \in (H \otimes H)^*$  satisfying  $r * \bar{r} = \varepsilon = \bar{r} * r$
- $\mu^{\circ p} = r * \mu * \bar{r}$

← such  $r$  is called a universal r-form

b) A quasi-cocommutative bialgebra  $H$  with universal  $r$ -form  $r$  is cobracted if

$$r(\mu \otimes \text{id}_H) = \tau_{13} * \tau_{23} \quad \text{and} \quad r(\text{id}_H \otimes \mu) = \tau_{13} * \tau_{12}$$

$$\text{where } \tau_{12} = r \otimes \varepsilon, \tau_{23} = \varepsilon \otimes r, \tau_{13} = (\varepsilon \otimes r)(\tau_{12} \otimes \text{id})$$

Here,  $*$  is the convolution (see Lecture 2). For example, the equality  $r * \bar{r} = \varepsilon$  means the following:

$$\forall x \otimes y \in H \otimes H \quad \sum_{(x)(y)} r(x' \otimes y') \cdot \bar{r}(x'' \otimes y'') = \varepsilon(x)\varepsilon(y).$$

Remark: Let's explain how these properties arise naturally by dualizing the  $R$  for  $R \in H \otimes H$  from p.1. To this end, we will draw diagrams and revert arrows.

(1)  $R$ -invertible means  $\exists R^{-1} \in H \otimes H$  s.t.  $R \cdot R^{-1} = \text{id}_{H \otimes H} = R^{-1} \cdot R$ .

Viewing  $R$  as a linear map  $\mathbb{K} \xrightarrow{R} H \otimes H$ ,  $R^{-1}$  as a linear map  $\mathbb{K} \xrightarrow{R^{-1}} H \otimes H$

that means

$$\mathbb{K} = \mathbb{K} \otimes \mathbb{K} \xrightarrow{R \otimes R^{-1}} (H \otimes H) \otimes (H \otimes H) \xrightarrow[\begin{smallmatrix} 1 \mapsto 1 \otimes 1 \\ \text{id} \otimes \text{id} \end{smallmatrix}]{} (H \otimes H) \otimes (H \otimes H) \xrightarrow{\mu \otimes \mu} H \otimes H$$

Revert the arrows, we get (by swapping  $\mu \leftrightarrow \Delta$ ,  $\varepsilon \leftrightarrow r$ ,  $R \leftrightarrow \bar{r}$ ,  $R^{-1} \leftrightarrow \bar{r}$ )

$$\mathbb{K} = \mathbb{K} \otimes \mathbb{K} \xleftarrow{r \otimes \bar{r}} (H \otimes H) \otimes (H \otimes H) \xleftarrow[\begin{smallmatrix} \text{id} \otimes \bar{r} \otimes \text{id} \\ \varepsilon_H \otimes \varepsilon_H \end{smallmatrix}]{} (H \otimes H) \otimes (H \otimes H) \xleftarrow{\Delta \otimes \Delta} H \otimes H$$

which precisely encodes  $r * \bar{r} = \varepsilon$ .

Swapping the order  $R \leftrightarrow R^{-1}$  in the first diagram will yield a swapped  $\tau \leftrightarrow \bar{\tau}$   $\Delta \leftrightarrow \varepsilon$  diagram encoding  $\bar{r} * r = \varepsilon$ .

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(Continuation of Runk)

(2) Let's now interpret the property  $\Delta^{\text{op}} = R \Delta R^{-1}$  via diagrams:

$$\begin{array}{ccccc}
 & H & \xrightarrow{\Delta} & H \otimes H & \xrightarrow{\tau} H \otimes H \xleftarrow{\mu^2 \otimes \mu^2} \\
 & \parallel & & & \\
 k \otimes H \otimes k & \longrightarrow & (H \otimes H) \otimes (H \otimes H) \otimes (H \otimes H) & \xrightarrow{\tau_{34} \tau_{45} \tau_{23}} & (H \otimes H \otimes H) \otimes (H \otimes H \otimes H) \\
 & & R \otimes \Delta \otimes R^{-1} & &
 \end{array}$$

Dualizing this diagram by reverting all arrows (swapping  $\Delta \mapsto \mu$ ,  $\mu \mapsto \Delta$ ) get:

$$\begin{array}{ccccc}
 & H & \xleftarrow{\mu^{\text{op}}} & H \otimes H & \xleftarrow{\tau} H \otimes H \\
 & \parallel & & & \\
 k \otimes H \otimes k & \xleftarrow{\tau \otimes \mu \otimes \tau} & (H \otimes H) \otimes (H \otimes H) \otimes (H \otimes H) & \xleftarrow{\tau_{23} \tau_{45} \tau_{34}} & (H \otimes H \otimes H) \otimes (H \otimes H \otimes H) \\
 & & R \otimes \Delta \otimes R^{-1} & &
 \end{array}$$

which precisely encodes  $\mu^{\text{op}} = \tau * \mu * \bar{\tau}$

(3) Let's now interpret  $(\Delta \otimes \text{id})(R) = R_{13} R_{23}$  via:

$$\begin{array}{ccccc}
 k & \xrightarrow{R} & H \otimes H & \xrightarrow{\Delta \otimes \text{id}} & H \otimes H \otimes H \\
 & \searrow R_{13} \circ R_{23} & (H \otimes H \otimes H) \otimes (H \otimes H \otimes H) & \xrightarrow{\text{id} \otimes H \otimes \text{id}} & H \otimes H \otimes H \\
 & & R_{13} \cdot R_{23} & &
 \end{array}$$

Dualize this diagram, we get:

$$\begin{array}{ccccc}
 k & \xleftarrow{\tau} & H \otimes H & \xleftarrow{\text{id} \otimes \text{id}} & H \otimes H \otimes H \\
 & \nearrow \tau_{13} * \tau_{23} & (H \otimes H \otimes H) \otimes (H \otimes H \otimes H) & \xleftarrow{H \otimes \text{id} \otimes H} & H \otimes H \otimes H \\
 & & \tau_{13} \cdot \tau_{23} & &
 \end{array}$$

Likewise, one can easily see that  $\tau_{13} * \tau_{12} = \tau(\text{id}_H \otimes \mu)$  is indeed dual to  $(\text{id} \otimes \Delta)(R) = R_{13} R_{12}$ .

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Similar to how a braiding on  $H$  was giving a braiding on the category  $H\text{-Mod}$ , we shall now see that a cobrading on  $H$  gives a braiding on the category  $H\text{-Comod}$ . The best way to see the construction is again by dualizing diagrams.

$$c_{v,w}^R: V \otimes W = k \otimes V \otimes W \xrightarrow{R \otimes \text{id}} (H \otimes H) \otimes V \otimes W \xrightarrow{\text{id} \otimes R \otimes \text{id}} H \otimes V \otimes H \otimes W \xrightarrow{\mu_H \otimes \mu_W} V \otimes W \xrightarrow{\tau} W \otimes V$$

$\downarrow$  dualize

$$V \otimes W = k \otimes V \otimes W \xleftarrow{\text{id} \otimes \text{id}} (H \otimes H) \otimes V \otimes W \xleftarrow{\text{id} \otimes \tau \otimes \text{id}} H \otimes V \otimes H \otimes W \xleftarrow{\Delta_V \otimes \Delta_W} V \otimes W \xleftarrow{\tau} W \otimes V$$

$\downarrow$  relabel  $V \rightsquigarrow W$ ,  $W \rightsquigarrow V$

$$c_{v,w}^L: V \otimes W \xrightarrow{\tau} W \otimes V \xrightarrow{\Delta_W \otimes \Delta_V} (H \otimes W) \otimes (H \otimes V) \xrightarrow{\tau \circ \text{id}} H \otimes H \otimes W \otimes V \xrightarrow{\text{id} \otimes \text{id}} W \otimes V$$

Using the Sweedler's notation  $\Delta_V(v) = \sum_{(v)} v_H \otimes v_V$ ,  $\Delta_W(w) = \sum_{(w)} w_H \otimes w_W$

(see Lecture 3), we can write an explicit algebraic formula for  $c_{v,w}^L$ :

$$c_{v,w}^L(v \otimes w) = \sum_{(v)(w)} c(w_H \otimes v_H) \cdot w_W \otimes v_V$$

The following result is dual to Proposition 2

Proposition 4: The above maps  $c_{v,w}^L$  satisfy:

- a)  $c_{v,w}^L$  - isomorphism of  $H\text{-comodules}$
- b)  $\forall U, V, W \in H\text{-Comod}$

$$c_{u \otimes v, w}^L = (c_{u,w}^L \otimes \text{id}_v)(\text{id}_u \otimes c_{v,w}^L)$$

$$c_{u,v \otimes w}^L = (\text{id}_v \otimes c_{u,w}^L)(c_{u,v}^L \otimes \text{id}_w)$$

- c)  $\forall U, V, W \in H\text{-Comod}$

$$(c_{v,w}^L \otimes \text{id}_u)(\text{id}_v \otimes c_{u,w}^L)(c_{u,v}^L \otimes \text{id}_w) = (\text{id}_w \otimes c_{u,v}^L)(c_{u,w}^L \otimes \text{id}_v)(\text{id}_u \otimes c_{v,w}^L)$$

! In other words, similarly to Prop 2 setup, the above result yields that any cobraided bialgebra  $H$  allows for a natural braiding on the  $\otimes$ -category of left  $H$ -comodules

Exercise: Prove Prop. 4

Hint: The inverse map  $\bar{c}_{v,w}^L: W \otimes V \rightarrow V \otimes W$  is given by  $\bar{c}_{v,w}^L(w \otimes v) = \sum_{(v)(w)} \bar{c}(w_H \otimes v_H) \cdot v_V \otimes w_W$

We end with the famous "FRT (=Faddeev-Reshetikhin-Takhtajan) construction".

Following [Kassel, VII.6] we split it into 2 theorems:

Theorem 1: Let  $V$  be a fin. dim. vector space and  $c \in \text{End}(V \otimes V)$ . Then, there is a (unique up to isomorphism) bialgebra  $A(c)$  with a linear map  $\Delta_V: V \rightarrow A(c) \otimes V$  s.t.

1)  $\Delta_V$  is a comodule map

2)  $c: V \otimes V \otimes V$  is a comodule morphism

3)  $A(c)$  is "universal" in the sense that given any bialgebra  $A'$  coacting on  $V$

via  $\Delta'_V: V \rightarrow A' \otimes V$  s.t.  $c$ - $A'$ -comodule morphism  $\exists! A(c) \xrightarrow{\cong} A'$  s.t.

$$\begin{array}{ccc} V & \xrightarrow{\Delta_V} & A(c) \otimes V \\ & & \searrow \xrightarrow{\text{coid}} \\ & & A' \otimes V \\ & & \Delta'_V \end{array}$$

Theorem 2: In the above setup if  $c$  is an R-matrix (i.e. invertible and satisfies Yang-Baxter equation)

then  $\exists!$  form  $r \in (A(c) \otimes A(c))^*$  making  $A(c)$  into a cobraided bialgebra s.t.

$$C_{V,V}^r = c.$$

Let's sketch the proofs of these theorems

Proof of Thm 1

• First, we define  $A(c)$  as an algebra by generators & relations. To this end, pick a basis  $v_i v_j$  of  $V$  and write  $c$  in the basis  $v_i \otimes v_j$   $\forall i,j=1 \dots N$ :

$$c(v_i \otimes v_j) = \sum_{m,n=1}^N C_{ij}^{mn} v_m \otimes v_n, \text{ so that } (C_{ij}^{mn}) - \text{matrix coeffs}$$

Define

$$A(c) = \mathbb{K}\langle T_i^j \mid 1 \leq i,j \leq N \rangle / I(c)$$

where  $I(c)$  is the 2-sided ideal gen'd by

$$A_{ij}^{mn} := \sum_{k,l=1}^N C_{ij}^{kl} T_k^m T_l^n - \sum_{k,l=1}^N C_{ki}^{ml} T_k^l T_j^m$$

← these will come naturally when verifying 2)

• To define the bialgebra structure it's convenient to form an  $N \times N$  matrix

$$T = (T_i^j)_{i,j=1}^N \quad \text{with } T_i^j \text{ being the } (i,j)^{\text{th}} \text{-entry.}$$

Lemma 1:  $\exists!$  bialgebra structure on  $A(c)$  with  $\Delta(T) = T \otimes T$ ,  $\epsilon(T) = I_N$

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On the level of generators we get:  $\Delta(T_i^j) = \sum_k T_i^k \otimes T_k^j, \quad \varepsilon(T_i^j) = \delta_{ij}$

- The above  $\varepsilon$ -las clearly define a bialg. structure on the free alg. gen'd by  $T_i^j$ . Thus, we need to check  $I(c)$  is a coideal:

$$1) \varepsilon(A_{ij}^{mn}) = 0 \text{ clearly}$$

$$2) \varepsilon(A_{ij}^{mn}) = \sum_{p,q} A_{ij}^{pq} \otimes T_p^m T_q^n + \sum_{p,q} T_i^p T_j^q \otimes A_{pq}^{mn}$$

- Coaction of  $A(c)$  on  $V$  is best written in matrix form

$$\boxed{\Delta_v \begin{pmatrix} v_1 \\ \vdots \\ v_n \end{pmatrix} = T \otimes \begin{pmatrix} v_1 \\ \vdots \\ v_n \end{pmatrix}}$$

i.e.  $\boxed{\Delta_v(v_i) = \sum_{j=1}^n T_i^j \otimes v_j}$ . The matrix algebraic form easily imply its coaction.

- The key part is to check  $c$  becomes a comodule morphism

Note, by definition, that:  $\boxed{\Delta_{v \otimes v}(v_i \otimes v_j) = \sum_{k,l} T_i^k T_j^l \otimes v_k \otimes v_l}$

$c$ -comodule morphism

$\Updownarrow$

$$\Delta_{v \otimes v}(c(v_i \otimes v_j)) - (\text{id}_{A(c)} \otimes c)(\Delta_{v \otimes v}(v_i \otimes v_j)) = 0 \quad \forall i,j$$

$\Downarrow$

$$\sum_{m,n} \left( \underbrace{\sum_{k,l} (c_{ij}^{kl} T_k^m T_l^n - c_{kj}^{mn} T_i^k T_j^l)}_{A_{ij}^{mn}} \right) \otimes v_m \otimes v_n = 0 \quad \forall i,j$$

$\Updownarrow$

$$A_{ij}^{mn} = 0 \quad \forall i,j, m,n \quad \leftarrow \text{which is by def'n of } A(c).$$

- Universal property is now immediate from the previous bullet!

Sketch of Proof of Theorem 2

- Let's first deduce the values of  $\tau$  on generators. To this end, we apply  $\stackrel{\tau}{c}_{v,w}(v \otimes w) = \sum \tau(w_h \otimes v_h) w_h \otimes v_h$  in the context  $v = v_j, w = v_i$
- $c_{v,w}^r(v \otimes w) = \sum_{m,n} c_{ji}^{mn} v_m \otimes w_n$

$w = V$   
 $H = A(c)$   
 $v = v_j, w = v_i$
- $\Rightarrow c(v_j \otimes v_i) = \sum_{m,n} c_{ji}^{mn} v_m \otimes v_n$

$\left. \begin{array}{l} \\ \end{array} \right\} \Rightarrow \boxed{c(T_i^m \otimes T_j^n) = c_{ji}^{mn}}$
- $c_{v,v}^r(v_j \otimes v_i) = \sum_{m,n} c(T_i^m, T_j^n) v_m \otimes v_n$

$\left. \begin{array}{l} \\ \end{array} \right\} \Rightarrow \boxed{c(T_i^m \otimes T_j^n) = c_{ji}^{mn}}$
- Cobracket properties recover  $\tau$  on the entire algebra uniquely by knowing the values on generators. Thus, it suffices to show such  $\tau$  exists

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(Continuation)

Exercise: Verify  $\tau(T_p^q \otimes A_{ij}^{mn}) = \tau(A_{ij}^{mn} \otimes T_p^q) \quad \forall i,j,p,q$  is a consequence of VB one where both  $\tau(\dots)$  are computed using cobracket property and values on generators.

Once such  $\tau$  is now constructed, verifying its cobracket is straightforward

Exercise: Verify  $\tau c (A(c) \otimes A(c))^*$  is a cobracket

We note that our construction above is very similar to that of  $M_q(2)$  and its coaction on  $\mathbb{K}^2$  from Lecture 6. In fact,  $M_q(2)$  can be recovered that way:

Exercise: Let  $V$  be a 2-dim v.space with a basis  $\{v_0, v_1\}$ , then  $V \otimes V$  is 4-dim with a basis  $\{v_0 \otimes v_0, v_0 \otimes v_1, v_1 \otimes v_0, v_1 \otimes v_1\}$ . Recall that

$$c = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & q^{-1} & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \in \text{End}(V \otimes V)$$

is a solution of VB eqn (which was obtained in lectures 14-15 by viewing  $V$  as  $L(1,+) \in U_q(\mathfrak{sl}_2)$ -mod, so that  $c$  is actually  $U_q(\mathfrak{sl}_2)$ -mod morphism)

Theorem:  $A(c) \cong M_q(2)$  (compare to [Homework 6, Problem 2])