

Primal and Dual Problems

(P)

$$\begin{aligned} \max \quad & \zeta(x) = c^T x \\ \text{subject to} \quad & Ax \leq b \\ & x \geq 0 \end{aligned}$$

(D)

$$\begin{aligned} \min \quad & \zeta(y) = b^T y \\ \text{subject to} \quad & A^T y \geq c \\ & y \geq 0 \end{aligned}$$



$$\begin{aligned} \max \quad & -\zeta(y) = -b^T y \\ \text{subject to} \quad & (-A^T) y \leq -c \\ & y \geq 0 \end{aligned}$$

Primal and Dual Problems

$$(P) \quad \begin{array}{l} \max \quad \zeta(x) = c^T x \\ \text{subject to} \quad Ax \leq b \\ \quad \quad \quad x \geq 0 \end{array} \quad \sim \quad \sum_j c_j x_j$$

$$\sum_i y_i \left(\sum_j a_{ij} x_j \leq b_i \right), \quad y_i \geq 0$$

$$\sum_j \underbrace{\left(\sum_i a_{ij} y_i \right)}_{c_j \leq} x_j \leq \sum_i b_i y_i$$

Thm 5.1 Weak Duality

$$\zeta(x) = \sum_j c_j x_j \leq \sum_i b_i y_i = \zeta(y)$$

Matrix Representations

$$(P) \max \hat{J}(X) = \hat{J}_0 + C^T X$$

$$W = b - AX \geq 0$$

$$X \geq 0$$

$$\begin{pmatrix} w_1 \\ w_2 \\ \vdots \\ w_m \end{pmatrix}$$

$$(D) \max -\hat{Z}(Y) = -\hat{Z}_0 - b^T Y$$

$$Z = -C + A^T Y \geq 0$$

$$Y \geq 0$$

$$\begin{pmatrix} z_1 \\ z_2 \\ \vdots \\ z_n \end{pmatrix}$$

\hat{J}_0	C_1	C_2	...	C_n
b_1	$-a_{11}$	$-a_{12}$...	$-a_{1n}$
b_2	$-a_{21}$	$-a_{22}$...	$-a_{2n}$
		\vdots		\vdots
b_m	$-a_{m1}$	$-a_{m2}$...	$-a_{mn}$

$-\hat{Z}_0$	$-b_1$	$-b_2$...	$-b_m$
$-C_1$	a_{11}	a_{21}	...	a_{m1}
$-C_2$	a_{12}	a_{22}	...	a_{m2}
	\vdots			
$-C_n$	a_{1n}	a_{2n}	...	a_{mn}

Matrix Representations

$$(P) \max \hat{J}(x) = \hat{J}_0 + C^T X$$

$$W = b - AX \geq 0$$

$$X \geq 0$$

$$\begin{pmatrix} w_1 \\ w_2 \\ \vdots \\ w_m \end{pmatrix}$$

$$(D) \max -\hat{\xi}(y) = -\hat{\xi}_0 - b^T y$$

$$Z = -C + A^T y \geq 0$$

$$y \geq 0$$

$$\begin{pmatrix} z_1 \\ z_2 \\ \vdots \\ z_n \end{pmatrix}$$

\hat{J}_0	C^T
b	$-A$ negative transpose
$-\hat{\xi}_0$	$-b^T$
$-C$	A^T

Complementary Variables

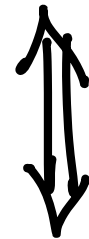
$(x_1, x_2, \dots, x_n, w_1, w_2, \dots, w_m)$

$y_i \sim w_i$

$x_j \sim z_j$

$(y_1, y_2, \dots, y_m, z_1, z_2, \dots, z_n)$

(P) x_j enter (leaves), w_i enter (leaves)



(D) z_j leaves (enters), y_i leaves (enters)

Neg. Transpose Property is Preserved in each Simplex Step

[V] p.64

(P)	$\zeta = +4x_1 + 1x_2 + 3x_3$
	$w_1 = 1 - x_1 - 4x_2$
	$w_2 = 3 - 3x_1 + x_2 - x_3$
- $\xi = -1y_1 - 3y_2$	
(D)	$z_1 = -4 + y_1 + 3y_2$
	$z_2 = -1 + 4y_1 - y_2$
	$z_3 = -3 + y_2$

x_3 enters, w_2 leaves
(greatest increase rule)

neg. transpose

z_3 leaves, y_2 enters

Neg. Transpose Property is Preserved in each Simplex Step

[V] p.64

$$\begin{array}{r} \zeta = 9 - 5x_1 + 4x_2 - 3w_2 \\ \hline (P)_2 \quad w_1 = 1 - x_1 - 4x_2 + \\ x_3 = 3 - 3x_1 + x_2 - w_2 \end{array}$$

$$\begin{array}{r} -\xi = -9 - 1y_1 - 3z_3 \\ \hline (D)_2 \quad z_1 = 5 + y_1 + 3z_3 \\ z_2 = -4 + 4y_1 - z_3 \\ y_2 = 3 + z_3 \end{array}$$

x_2 enters, w_1 leaves

neg. transpose

z_2 leaves, y_1 enters

Neg. Transpose Property is Preserved in each Simplex Step

[V] p.64

(P)

$\zeta = 10$	-6	x_1	-1	w_1	-3	w_2
$x_2 =$	0.25	-0.25	x_1	-0.25	w_1	
$x_3 =$	3.25	-3.25	x_1	-0.25	w_1	w_2

(D)

$-\xi = -10$	-0.25	z_2	-3.25	z_3
$z_1 =$	6	$+0.25$	z_2	$+3.25$
$y_1 =$	1	$+0.25$	z_2	$+0.25$
$y_2 =$	3		$+1$	z_3

Optimal for (P)

neg. transpose

Optimal for (D)

Optimal values for y_1, y_2 ($z_2, z_3 = 0, y_1 = 1, y_2 = 3, z_1 = 0$)

Neg. Transpose Property is Preserved in each Simplex Step

at Optimal:

(P)

S^*	C_1^*	C_2^*	...	C_n^*
b_1^*	$-a_{11}^*$	$-a_{12}^*$...	$-a_{1n}^*$
b_2^*				
\vdots				
b_m^*	$-a_{m1}^*$	$-a_{m2}^*$...	$-a_{mn}^*$

$C_j^* \leq 0$

$b_i^* \geq 0$

Neg. Transpose Property is Preserved
in each Simplex Step

at Optimal:

(D)

Strong
Duality

$-c_j^* \geq 0$

$-b_i^* \leq 0$

$-z^*$	$-b_1^*$	$-b_2^*$...	$-b_m^*$
$-c_1^*$	a_{11}^*	a_{21}^*	...	a_{m1}^*
$-c_2^*$				
...			a_{ji}^*	
$-c_n^*$	a_{1n}^*	a_{2n}^*	...	a_{mn}^*

Strong Duality Theorem

If (P) has an optimal solution (x_j^*)
then (D) also has an optimal solution (y_i^*)
and

$$(\sum^*) \sum_j c_j x_j^* = \sum_i b_i y_i^* (= \sum^*)$$

is. duality gap = 0

(The proof gives a formula for y_i^*)

Strong Duality Theorem

Pf (\mathcal{N} : non-basic vars)
(\mathcal{B} : basic vars)

$$\textcircled{1} \quad \mathcal{J} = \sum_{\substack{j \\ \mathcal{N}}} \bar{c}_j x_j + \sum_{\substack{i \\ \mathcal{B}}} 0 w_i$$

(original formulation)

Strong Duality Theorem

~~Pf~~


② at (P)'s optimal

$$J = J^* + \sum_j \bar{c}_j \bar{x}_j + \sum_i 0 \bar{w}_i$$

$(x_1, x_2, \dots, x_n, w_1, \dots, w_m)$

$$= J^* + \sum_j \bar{c}_j^* \bar{x}_j + \sum_i \bar{d}_i^* \bar{w}_i$$

$(\bar{c}_j^* \leq 0)$ $(\bar{d}_i^* \leq 0)$

$$= J^* + \sum_j \bar{c}_j^* \bar{x}_j + \sum_i \bar{d}_i^* \left(b_i - \sum_j \bar{a}_{ij} \bar{x}_j \right)$$


Strong Duality Theorem

Pf ② at (P)'s optimal

$$J = J^* + \sum_j c_j^* x_j + \sum_i d_i^* \left(b_i - \sum_j a_{ij} x_j \right)$$

$$= \underline{\left(J^* + \sum_i d_i^* b_i \right)}$$

$$+ \sum_j \underline{\left(c_j^* - \sum_i a_{ij} d_i^* \right)} x_j$$

Strong Duality Theorem

Pf Compare ① and ②

Const term: $0 = J^* + \sum_i d_i^* b_i$

x_j : $C_j = C_j^* - \sum_i a_{ij} d_i^*$

Let $y_i^* = -d_i^* (\geq 0)$. Then

$$J^* = \sum_i b_i y_i^*$$

Strong Duality Theorem

Pf Compare ① and ②

Const term: $0 = J^* + \sum_i d_i^* b_i$

x_j : $C_j = C_j^* - \sum_i a_{ij} d_i^*$

C_j = $C_j^* + \sum_i a_{ij} y_i^* \leq \sum_i a_{ij} y_i^*$
($C_j^* \leq 0$)

Strong Duality Theorem

Pf

Hence there exists an y^* s.t.

$$J^* = J(x^*) = \sum(y^*)$$



This must be opt. for (D)

(By Weak Duality,

$$J(x^*) \leq \sum(y^*))$$

Four and Only Four Possibilities

(1) (P) has an optimal solution

\Rightarrow (D) also has an optimal solution

$$J^* = Z^*, \text{ duality gap} = 0$$

(2) (P) is feasible but unbounded

\Rightarrow (D) must be infeasible

$$J(x) \leq Z(y)$$

$\rightarrow +\infty$

$$J^* = Z^* = +\infty$$

(min empty set = $+\infty$)

Four and Only Four Possibilities

(3) (D) is feasible but unbounded

\implies (P) must be infeasible

$$J^*(x) \leq \sum^*(y) \\ \rightarrow -\infty$$

$$J^* = \sum^* = -\infty \\ (\text{max empty set} = -\infty)$$

(4) Both (P) and (D) are infeasible

$$\implies J^* = -\infty, \quad \sum^* = +\infty$$

$$(\text{duality gap} = +\infty)$$

Four and Only Four Possibilities

[V, p.69]

(P)

$$\text{maximize } 2x_1 - x_2$$

$$\text{subject to } x_1 - x_2 \leq 1 \quad \leftarrow$$

$$-x_1 + x_2 \leq -2 \quad \leftarrow$$

$$x_1, x_2 \geq 0.$$

inconsistent

(D)

$$\text{min } y_1 - 2y_2$$

$$\text{subject to } y_1 - y_2 \geq 2 \quad \leftarrow$$

$$-y_1 + y_2 \geq -1 \quad \leftarrow$$

$$y_1, y_2 \geq 0$$

inconsistent

Complementary Slackness

[P]

THEOREM 5.3. Suppose that $x = (x_1, x_2, \dots, x_n)$ is primal feasible and that $y = (y_1, y_2, \dots, y_m)$ is dual feasible. Let (w_1, w_2, \dots, w_m) denote the corresponding primal slack variables, and let (z_1, z_2, \dots, z_n) denote the corresponding dual slack variables. Then x and y are optimal for their respective problems if and only if

$$(5.7) \quad \begin{aligned} x_j z_j &= 0, & \text{for } j = 1, 2, \dots, n, \\ w_i y_i &= 0, & \text{for } i = 1, 2, \dots, m. \end{aligned}$$

(pairing of variables)

i.e. either x_j or $z_j = 0$ $j=1, \dots, n$
and either y_i or $w_i = 0$ $i=1, \dots, m$

Complementary Slackness

Pf
$$\sum_j c_j x_j \leq \sum_j \left(\sum_i a_{ij} y_i \right) x_j$$
$$= \sum_i \left(\sum_j a_{ij} x_j \right) y_i \leq \sum_i b_i y_i$$

$$\sum_j c_j x_j \stackrel{\textcircled{1}}{=} \sum_j \left(\sum_i a_{ij} y_i \right) x_j$$
$$= \sum_i \left(\sum_j a_{ij} x_j \right) y_i \stackrel{\textcircled{2}}{\leq} \sum_i b_i y_i$$

Complementary Slackness

Pf

$$\sum_j c_j x_j = \sum_j \left(\sum_i a_{ij} y_i \right) x_j$$

$$0 = \sum_j \left(\underbrace{\sum_i a_{ij} y_i - c_j}_{z_j} \right) x_j$$

$$= \sum_j z_j x_j \geq 0$$

$$\implies z_j x_j = 0$$

Complementary Slackness

$$\text{Pf } \sum_i \left(\sum_j a_{ij} x_j \right) y_i = \sum_i b_i y_i \quad (2)$$

$$0 = \sum_i \left(b_i - \underbrace{\sum_j a_{ij} x_j}_{w_i} \right) y_i$$

$$= \sum_i w_i y_i \geq 0$$

$$\implies w_i y_i = 0$$