

Accelerated Proximal Gradient Methods

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outline

1 Background

2 Acceleration Methods

3 Numerical Results

Problem Statement

$$\min_{x \in \mathcal{E}} f(x) + P(x)$$

Assumptions

- \mathcal{E} is a linear space and $\text{dom} P \neq \emptyset$
- f is continuously differentiable
- ∇f is L-Lipschitz
- P is proximal

Terminology

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- $D(x, y) = h(x) - h(y) - \langle \nabla h(y), x - y \rangle$
- Example: $h(x) = \frac{1}{2}\|x\|^2 \rightarrow D(x, y) = \frac{1}{2}\|x - y\|^2$

Basic Proximal Gradient

$$x_{k+1} = (I + \alpha \partial P)^{-1} (I - \nabla f) x_k$$

- $\mathcal{O}(1/k)$ convergence
- $\mathcal{O}(n)$ memory

Provable Results

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- Duality gap shrinks with bounds depending on choice of momentum term, θ_k (q^P is dual function).

$$0 \leq f^P(x_{k+1} - q^P(\bar{v}_k)) \leq \theta_k^2 L \max_{x \in \text{dom}P} D(x, z_0)$$

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- Use “momentum” to accelerate
- Momentum decreases over time to hone in
- Sometimes overshoots, creating small oscillations

Algorithm 1

$$\begin{aligned}y_k &= (1 - \theta_k)x_k + \theta_k z_k \\z_{k+1} &= \arg \min_{x \in X_k} \{ \ell_f(x; y_k) + \theta_k LD(x, z_k) \} \\ \hat{x}_{k+1} &= (1 - \theta_k)x_k + \theta_k z_{k+1}\end{aligned}$$

with the constraints

$$\frac{1 - \theta_{k+1}}{\theta_{k+1}^2} \leq \frac{1}{\theta_k^2}$$
$$\ell_f(x_{k+1}; y_k) + \frac{L}{2} \|x_{k+1} - y_k\|^2 \leq \ell_f(\hat{x}_{k+1}; y_k) + \frac{L}{2} \|\hat{x}_{k+1} - y_k\|^2$$

Vert²

Algorithm 1

$$\theta_k = (1/2)(\sqrt{\theta_{k-1}^4 + 4\theta_{k-1}^2} - \theta_{k-1}^2)$$

$$y_k = (1 - \theta_k)x_k + \theta_k z_k$$

$$z_{k+1} = (I + \frac{1}{\theta_k L} \partial P)^{-1}(z_k - \frac{1}{\theta_k L} \nabla f(y_k))$$

$$x_{k+1} = (1 - \theta_k)x_k + \theta_k z_{k+1}$$

- $\mathcal{O}(1/k^2)$ convergence
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Algorithm 1

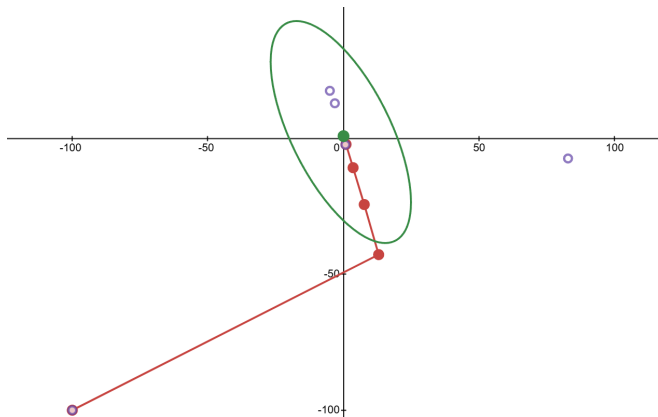


Figure: Red: x_k , Purple: z_k

Algorithm 2

$$y_k = x_k + \theta_k(\theta_{k-1}^{-1} - 1)(x_k - x_{k-1})$$
$$x_{k+1} = (I + \frac{1}{L}\partial P)^{-1}(I - \frac{1}{L}\nabla f)y_k$$

- This is the one we did in class, with $t_k = 1/\theta_k$.

Algorithm 3

$$y_k = (1 - \theta_k)x_k + \theta_k z_k$$

$$z_{k+1} = \arg \min_x \left\{ \sum_{i=0}^k \frac{\ell_f(x; y_i)}{\vartheta_i} + Lh(x) \right\}$$

$$x_{k+1} = (1 - \theta_k)x_k + \theta_k z_{k+1}$$

- Weighted sum of gradients from previous iterations
- Typically $\vartheta_k = \theta_k$, but technically can be relaxed.
- $\{\theta_k\}$ decreasing, so more recent terms weighted higher

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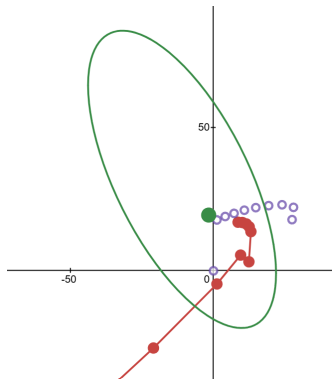


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Search Region Reduction

- $\forall w, x \in \text{dom}P$ such that $f(x) \leq \inf f + \epsilon$,
 $f^P(w) + \epsilon \geq f^P(x) \geq \ell_f(x; w)$ by convexity.

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- So convex combinations of these half spaces do too

$$X_k = \left\{ x : \sum_{i \in I_{k,j}} \alpha_{k,i} (\ell_f(x, w_{k,i}) - f^P(w_{k,i})) \leq \epsilon, j = 1, \dots, n_k \right\}$$

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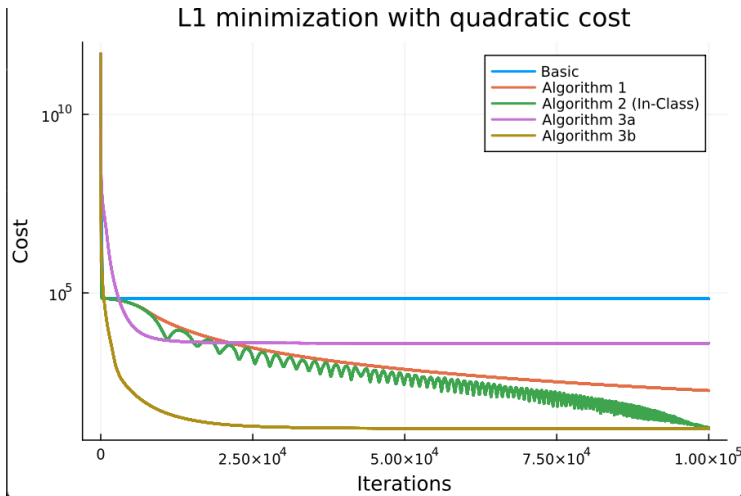
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- Half spaces relatively easy to search in, but still increases cost per iteration.

Not Strongly Convex



Strongly Convex

L1 minimization with quadratic cost

